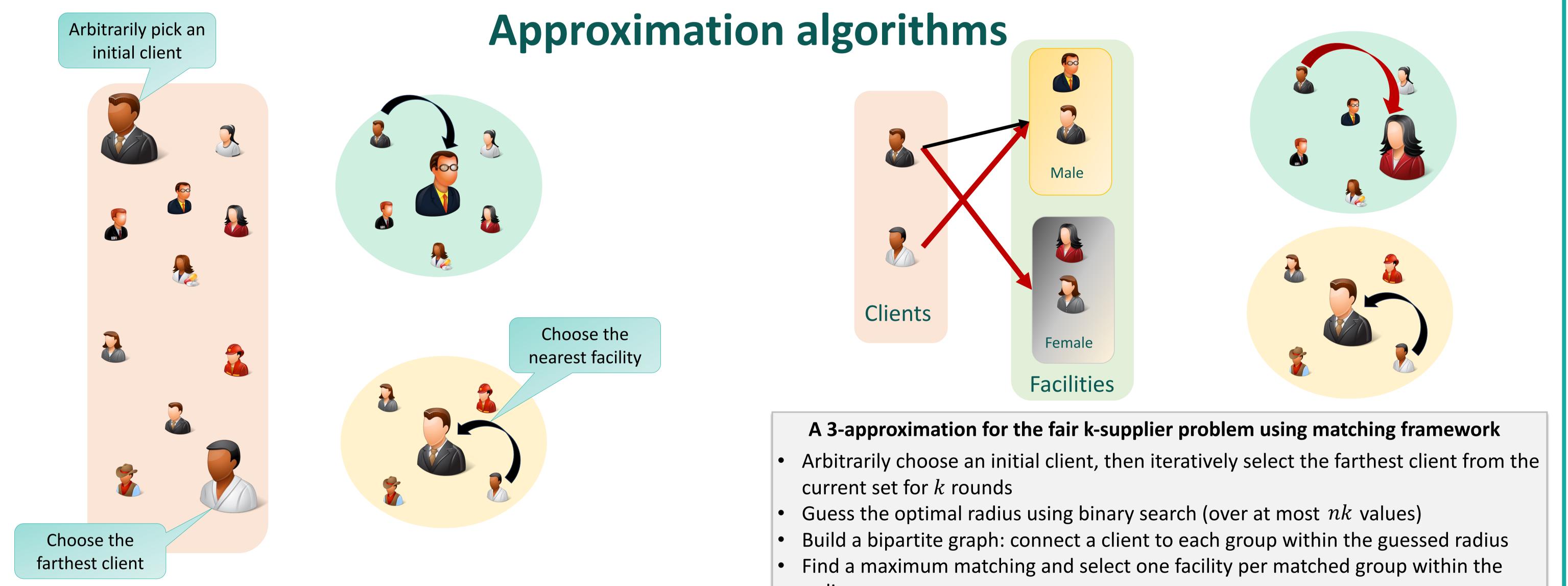


Cluster the data points by choosing k = 2 representative centers / facilities to minimize the maximum distance from any point to its closest representative

The goal remains to minimize the maximum distance, but with fairness constraints: each group must contribute a minimum (and a maximum) number of cluster centers

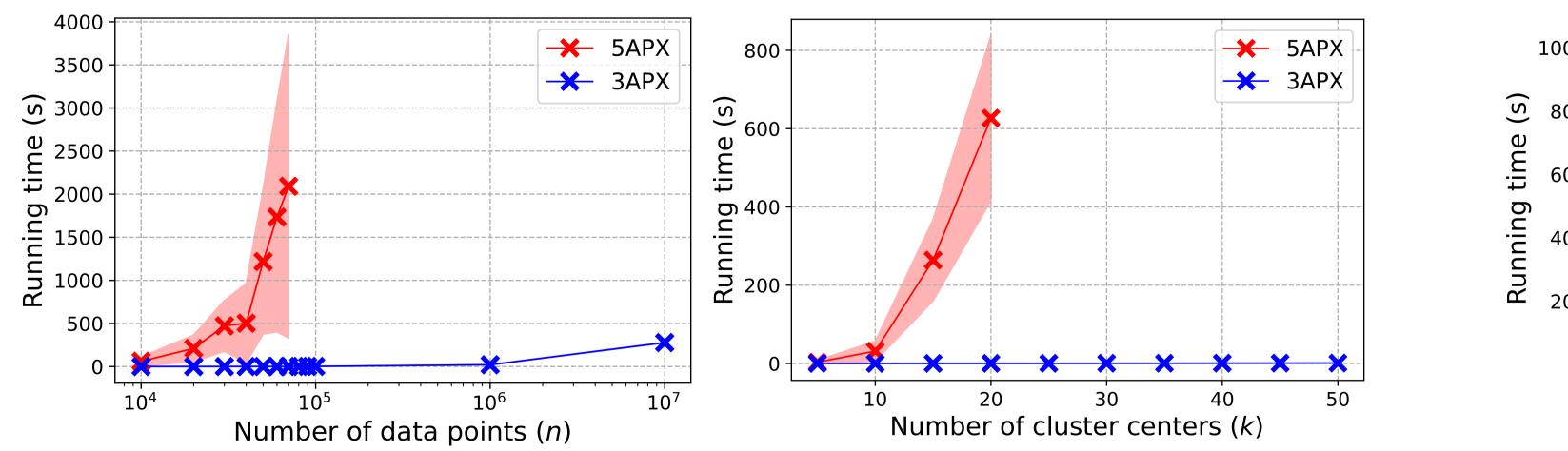


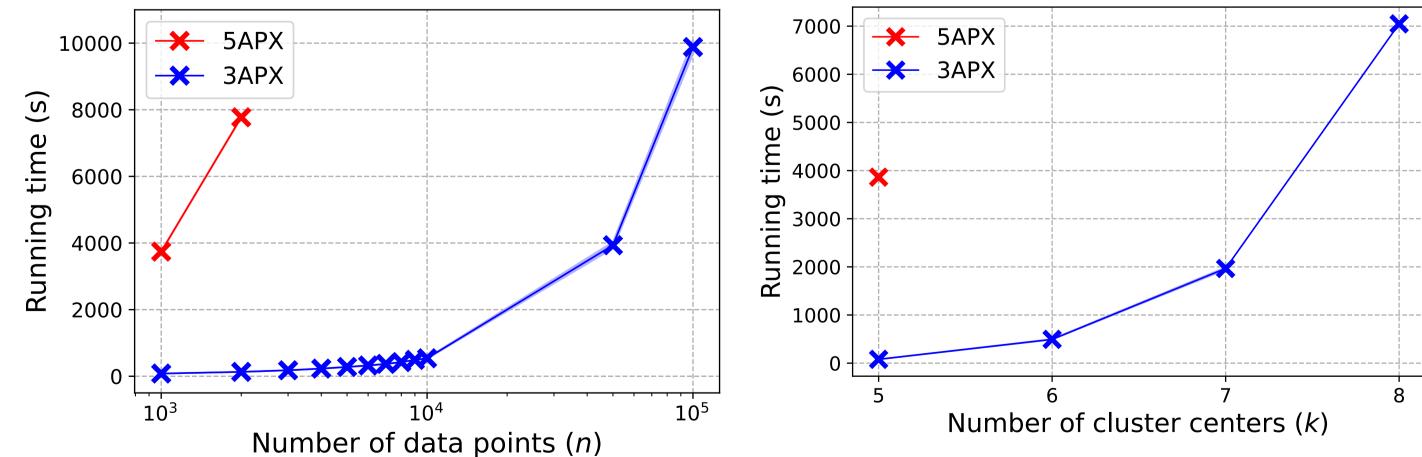
A greedy 3-approximation for the (unfair) k-supplier problem

- Arbitrarily choose an initial client, then iteratively select the farthest client \bullet from the current set for k rounds
- Choose the nearest facility to each selected client

- radius
- This yields a 3-approximation in $O((kn + k^2\sqrt{k})\log n\log k)$ time
- For intersecting groups, reduce the problem to multiple instances with disjoint groups, incurring an overhead of $O(2^{tk})$
- *n* is number of data points and *t* is number of groups

Experimental results





Salability of our 3-approximation algorithm (3-APX) vs. Chen et al.'s (TCS, 2024) 5-approximation (5-APX) for fair k-supplier with disjoint groups

- *Left plot*: Runtime scaling with increasing number of data points n, for k = 10 and 5 disjoint groups
- **Right plot:** Runtime scaling with increasing number of cluster centres k, for a fixed dataset of $n = 10\,000$ points
- *Fairness constraint:* select an equal number of cluster centres from each group

Salability of our 3-approximation algorithm (3-APX) vs. Thejaswi et al.'s (ArXiv, 2024) 5-approximation (5-APX) for fair k-supplier with intersecting groups

- *Left plot:* Runtime scaling with increasing number of data points n, for k = 5 and 5 intersecting groups. Fairness constraint: choose at least 2 cluster centres from each group
- **Right plot:** Runtime scaling with increasing number of cluster centres k, for a fixed dataset of n = 1000 points. Fairness constraint: choose $\frac{2\kappa}{r}$ cluster centres from each group